# **K-ordered Hamiltonian Graphs**

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#### Abstract

In this chapter, we review the following results proved in <sup>1, 2, 3, 4</sup>: (i) For  $k \ge 3$ , every (k+1)-Hamiltonian-connected graph is k-ordered. Determine f(k, n) if n is sufficiently large in terms of k. Let g(k,n) = n/2 + k/2 - 1 (ii) f(k,n) = g(k,n) if  $n \ge 1$  lk-3 (iii)  $f(k,n) \ge g(k,n)$  for any  $n \ge 2k$  and f(k,n) > g(k,n) if  $2k \le n \le 3k-6$  (iv) if G is a graph of order n with  $3\le k\le n/2$  and  $deg(u)+deg(v) \ge n+(3k-9)/2$  for every pair u, v of non-adjacent vertices of G, then G is k-ordered Hamiltonian.

Keywords: Adjacency, Connectedness, Hamilton, Vertices

## 1. Introduction

In this section, the concept of K-Hamiltonian connected graphs is introduced. The result is that, for  $k\geq 3$ , every (k+1)-hamiltonian connected graph is k-ordered is established. Also, necessary and sufficient conditions for a graph to be k-ordered are discussed. The notations used are recalled here. We write v1, (x1), v2, (x2), ..., vk-1, (xk-1), vk to indicate a v1-vk path that contains the k vertices v1, v2...vk, and possibly upto k-1 additional vertices, namely x1, x2, ...xk-1. Thus, for  $1 \leq j \leq k-1$ , (x1) indicates that some vertex, which we denote by xi, may (or) may not be present on the path. So, for each i ( $1 \leq i \leq k-1$ ), the vertices v1 and vi+1 are either adjacent or vi, vi+1 is a path. For example, u, v (x), y indicates the path u, v, y or the path u, v, x, y.

**PROPOSITION 1.1:** Let G be a Hamiltonian graph of order  $n \ge 3$ . If G is k-ordered,  $3 \le k \le n$ , then G is (k-1)-connected.

**PROOF:** Suppose, to the contract that G is not (k-1) connected. Then, there exists a cut set  $S = v_1, v_2, ..., v$  of G, where  $\leq k-2$ . Let  $v_{+1}$  and  $v_{+2}$  be vertices belonging to distict components of G-S. Consequently, every  $v_{+1}$ - $v_{+2}$  path in G contains vertices of S, so G contains no Hamiltonian cycle containing the vertices of the sequence  $v_1, v_2, ..., v$ ,  $v_{+1}, v_{+2}$  in this order. Hence G is not (+2) -ordered and so is not k-ordered.

**COROLLARY 1.2:** If G is a k-ordered Hamiltonian graph, then  $\delta(G) \ge k-1$ .

#### Sufficient conditions for k-ordered graphs

Many sufficient conditions have been given for Hamiltonian graphs. Two of the best known are by Dirac and Ore.

**THEOREM 1.3 (DIRAC):** Let G be a graph of order  $n \ge 3$ . If deg  $v \ge n/2$  for every vertex v of G, then G is Hamiltonian.

**THEOREM 1.4 (ORE):** Let G be a graph of order  $n \ge 3$ . If for every pair u, v of non-adjacent vertices of G, deg u+deg v  $\ge n$ , then G is hamiltonian.

**LEMMA 1.5:** Let G be a graph of order  $n \ge 3$  and let S:  $v_1, v_2, \dots, v_k$  be a sequence of k distinct vertices of G, where  $3 \le k \le n$ . If

deg u+deg v 
$$\ge$$
 n+2k-7

For every pair u, v of non adjacent vertices of G, then G contains a  $v_1 - v_k$  path of the type  $v_1$ ,  $(x_1)$ ,  $v_2$ , $(x_2)$ ,...  $(x_{k-1})$ ,  $v_k$  or a  $v_k$ - $v_{k-1}$  path of the type  $v_k$ ,  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ , ..., $(x_{k-2})$ ,  $v_{k-1}$ .

**PROOF:** Consider the vertices  $v_1$  and  $v_2$ . If  $v_1$  and  $v_2$  are adjacent, then G contains the path  $v_1$ ,vs:; otherwise, deg  $v_1 + \deg v_2 \ge n + 2k - 7 \ge n - 1$ , which implies the existence of a vertex x1 mutually adjacent to  $v_1$  and  $v_2$  and so  $v_1$ ,  $x_1$ ,  $v_2$  is a path in G. In any case, G contains a path  $v_1$ ,  $(x_1)$ ,  $v_2$ . Now, consider the vertex  $v_3$ . If  $v_2v_3 \in E(G)$ , then G contains the path  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $v_3$ . Suppose, then that  $v_2v_2 \in E(G)$ . Thus, deg  $v_2 + \deg v_2 \ge n + 2k - 7$  by hypothesis. Since G has

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order n, the vertices  $v_2$  and  $v_3$  are mutually adjacent to at least 2k-5 vertices. if k≥4, then G contains a vertex  $x_2$ distinct from  $v_1$  and  $x_1$  (if it exists) such that  $x_2$  is mutually adjacent to  $v_2$  and  $v_3$ . Hence G contains the path  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ ,  $v_3$ . Suppose, then that k=3 and that G contains to vertex distinct from v1 and x1 that is mutually adjacent to  $v_2$  and  $v_3$ . If  $v_3$  is adjacent to  $v_1$ , then  $v_3$ ,  $v_1$ , $(x_1)$ ,  $v_2$  is a path of G. If  $v_3$  is not adjacent to  $v_1$ , then  $v_2$  and  $v_3$  are mutually adjacent to 1 and to no other vertex, while  $v_2$  is adjacent to  $v_1$ . However, then  $v_1$ ,  $v_2$ ,  $x_1$ ,  $v_3$  is a path in G.

Proceeding inductively, assume that we have constructed a path  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ ,...,  $(x_{j-2})$ ,  $v_{j-1}$  in G. Suppose first that  $j \le k-1$ . It is to show that G contains a path  $v_1$ ,...,  $(x_1) v_2 (x_2) v_j$ . If  $v_{j-1}$  is adjacent to  $v_j$ . If  $v_{j-1}$  is adjacent to  $v_j$ , then G contains such a path. Suppose then that  $v_{j-1} v_j \in$ E(G). If G contains a vertex  $x_{j-1}$  distinct from the vertices on the path  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ ,...,  $(x_{j-2})$  such that  $x_{j-1}$  is mutually adjacent to  $v_{j-1}$  and  $v_j$ , then G contains a path of the designed type, otherwise  $v_{j-1}$  and  $v_j$  are mutually adjacent to at most  $2_{j-4}$  vertices and G contains at least  $n + 2k - 2_{j-1}$ vertices. Since  $k \ge j + 1$ , this produces a contradiction, and hence the desired claim is verified.

Finally, suppose that j = k, that is, assume that there is a path  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ ,...,  $(x_{k-2})$ ,  $v_{k-1}$  is constructed in G. It is to show that G contains either a path  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ ,...,  $v_{k-1}$ ,  $(x_{k-1})$ ,  $v_k$  or a path  $v_k$ ,  $v_1$ ,  $(x_1)$ ,  $v_2$ ,  $(x_2)$ ...,  $(x_{k-2})$ ,  $v_{k-1}$ . If either  $v_k v_{k-1} \in E(G)$  or  $v_k v_1 \in E(G)$ , then the desired result is obtained, thus assume that  $v_k$  is adjacent to neither  $v_{k-1}$  nor  $v_1$ . Also, if  $v_{k-1}$  and vk are mutually adjacent to some vertex other than the (at most)  $2_{k-5}$  vertices on the path  $(x_1)$ ,  $v_2$ ,  $(x_2)$ , ...,  $v_{k-2}$ ,  $(x_{k-2})$ , the proof is complete; so assume that this is not the case. Hence  $v_{k-1}$  and  $v_k$  are mutually adjacent to at most  $2_{k-5}$  vertices. Since G has order n, it follows that  $v_{k-1}$  and  $v_k$  are mutually adjacent to exactly  $2_{k-5}$  vertices (so all of the vertices  $x_1, x_2, ..., x_{k-2}$  exist), and every other vertex of G is adjacent to exactly one of  $v_{k-1}$ , and G contains the path  $v_1, x_1, v_2, x_2, ..., v_{k-2}, x_{k-2}, v_k$ , producing the desired path. Hence, the result.

## 2. References

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